

## Bruhat Preclosure

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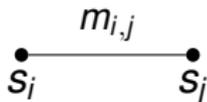
## Coxeter groups

- *Coxeter System*  $(W, S)$

$$W := \langle S \mid (s_i s_j)^{m_{i,j}} = e \text{ for } s_i, s_j \in S \rangle$$

$m_{i,j} \in \mathbb{N}^* \cup \{\infty\}$  and  $m_{i,j} = 1$  only if  $i = j$ .

- *Coxeter diagram*  $\Gamma_W$ : vertices  $S$  and edges labelled  $m_{i,j}$  when  $m_{i,j} > 3$ .



### Example

$$W_{B_3} = \langle s_1, s_2, s_3 \mid s_1^2 = s_2^2 = s_3^2 = (s_1 s_2)^4 = (s_2 s_3)^3 = (s_1 s_3)^2 = e \rangle$$



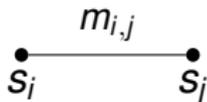
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### Example

$W_{A_n} = S_{n+1}$ , symmetric group.



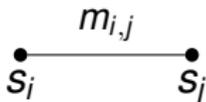
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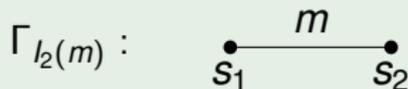
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## Example

$W_{I_2(m)} = \mathcal{D}(m)$ , dihedral group of order  $2m$ .



## Coxeter groups

■ Coxeter System  $(W, S)$ 

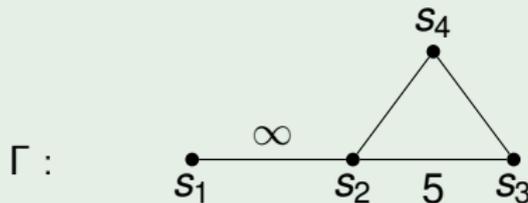
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■ Coxeter diagram  $\Gamma_W$ : vertices  $S$  and edges labelled  $m_{i,j}$  when  $m_{i,j} > 3$ .

## Example

$W$ , an infinite Coxeter group.



## Words

$(W, S)$  a Coxeter system.

- $w = s_1 \dots s_n \in W$  for  $s_i \in S$ .
- $w$  has *length*  $n$ ,  $\ell(w) = n$ , if  $n$  is minimal.
- If  $\ell(w) = n$  then  $w = s_1 \dots s_n$  is a *reduced expression*.

### Example

$$\Gamma_{A_2} : \begin{array}{c} s \quad t \\ \bullet \text{---} \bullet \end{array}$$

Then  $\ell(stst) = 2$

$$stst = tstt = ts$$

since

$$(st)^3 = e \Rightarrow sts = tst.$$

$ts$  is a reduced expression of  $stst$ .

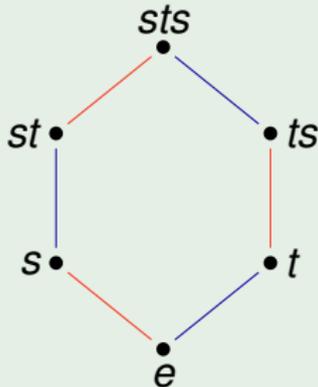
## Weak order

- *(right) weak order* is the prefix order on elements:

$$u \leq_R v \iff l(v) = l(u) + l(u^{-1}v)$$

### Example

$\Gamma_{A_2} : s \text{ --- } t$



## Inversion sets

$T = \cup wSw^{-1}$  is the set of *reflections*.

For  $w \in W$ , (*left*) *inversion set*:

$$N(w) = \{t \in T \mid \ell(tw) < \ell(w)\}$$

## Example

$$\Gamma_{A_2} : \bullet \xrightarrow{s} \bullet \xrightarrow{t} \bullet$$

Recall  $\ell(ts) = 2$ . Then

$$N(ts) = \{t, tst\}$$

since

$$\begin{aligned} (t)(ts) &= \mathbf{t}ts = s \\ (tst)(ts) &= \mathbf{t}stts = t \end{aligned}$$

## Generating inversion sets

For  $w = s_1 \dots s_n$  a reduced expression:

$$N(w) = \{s_1, s_1 s_2 s_1, s_1 s_2 s_3 s_2 s_1, \dots, s_1 s_2 \dots s_n \dots s_2 s_1\}$$

### Example

$$\Gamma_{A_2} : \bullet \xrightarrow{s} \bullet \xrightarrow{t}$$

$$N(ts) = \{t, tst\}$$

### Example

$$\Gamma_{A_3} : \bullet \xrightarrow{r} \bullet \xrightarrow{s} \bullet \xrightarrow{t}$$

$$N(rstsr) = \{r, rsr, rstsr, rststsr = t, rstsrstsr = tst\}$$

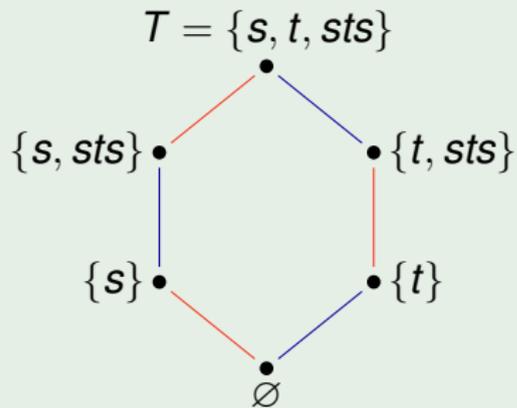
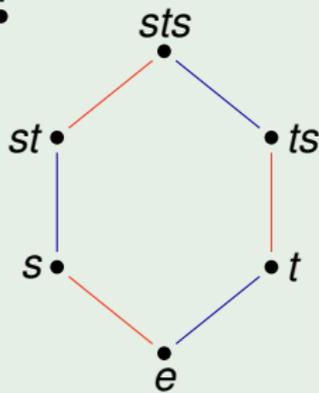
## Weak order - Round 2

- The *(right) weak order* is the prefix order on elements:

$$u \leq_R v \iff \ell(v) = \ell(u) + \ell(u^{-1}v) \iff \mathbf{N}(u) \subseteq \mathbf{N}(v)$$

### Example

$$\Gamma_{A_2} : \begin{array}{c} s \quad t \\ \bullet \text{---} \bullet \end{array}$$



## (Semi)lattices

$(W, S)$  a Coxeter system and  $u, v \in W$ .

- **Meet** (greatest lower bound) -  $u \wedge_R v$
- **Join** (least upper bound) -  $u \vee_R v$
- **Meet-semilattice** every two elements has a meet.
- **Join-semilattice** every two elements has a join.
- **Lattice** every two elements has a meet and join.

Theorem (Björner [1984])

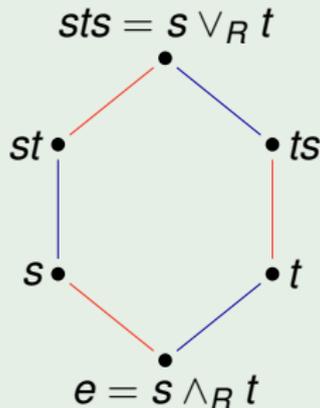
*The weak order is a meet-lattice. If  $W$  is finite, it's a lattice.*

## Lattice

- *Meet* (greatest lower bound) -  $u \wedge_R v$
- *Join* (least upper bound) -  $u \vee_R v$
- *Meet-semilattice* every two elements has a meet.
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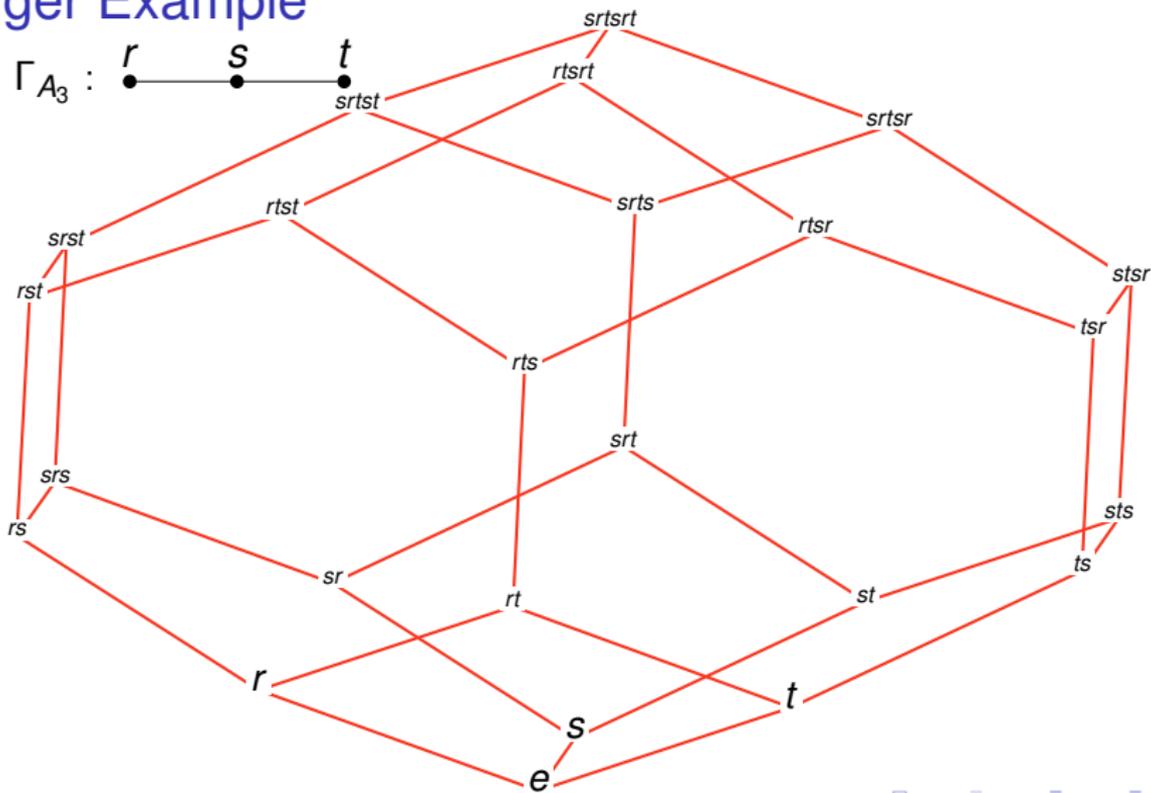
### Example

$$\Gamma_{A_2} : \begin{array}{c} s \quad t \\ \bullet \text{---} \bullet \end{array}$$



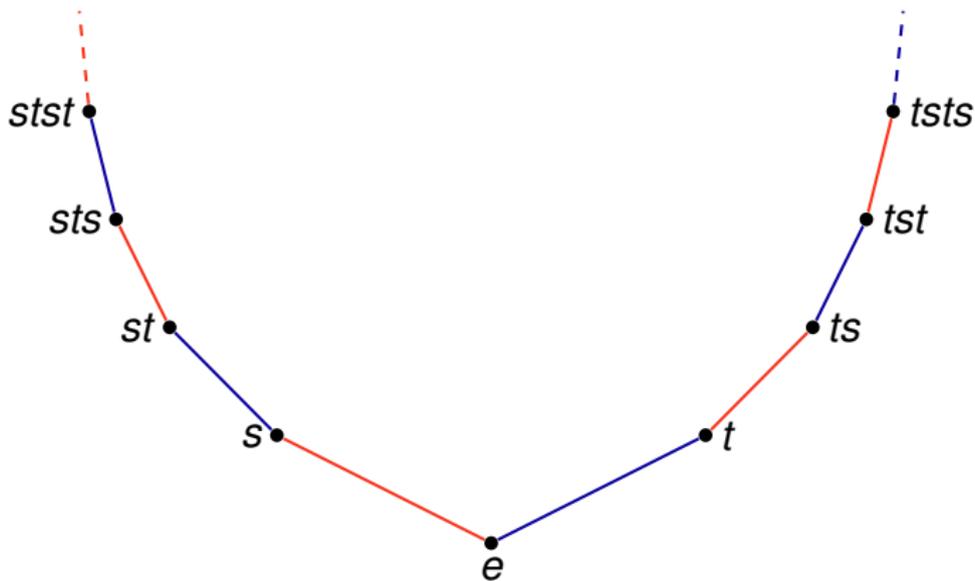
## Larger Example

$\Gamma_{A_3} : \bullet \xrightarrow{r} \bullet \xrightarrow{s} \bullet \xrightarrow{t}$



## Infinite example

$$\Gamma_{\tilde{A}_1} : \begin{array}{c} s \quad \infty \quad t \\ \bullet \text{---} \bullet \end{array}$$

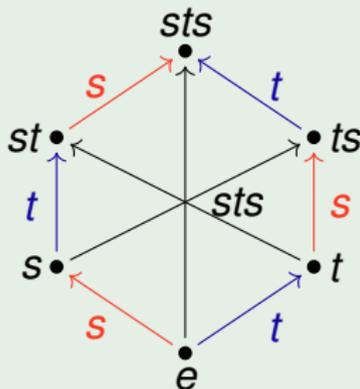


## Bruhat graph $\Omega$

- Vertices: Elements of  $W$
- Edges:  $(u, v)$  if
  - $v = ut$  for some  $t \in T$
  - $\ell(u) \leq \ell(v)$ .

### Example

$$\Gamma_{A_2} : \bullet \xrightarrow{s} \bullet \xrightarrow{t} \bullet$$

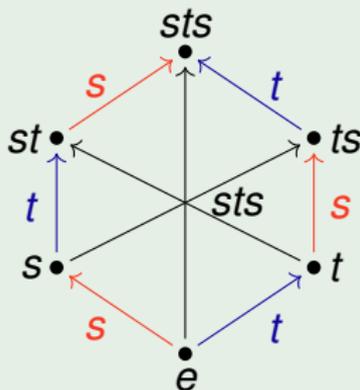


## A-Path

For  $A \subseteq T$ , *A-path* a directed path in  $\Omega$  following only edges labelled by elements in  $A$ .

### Example

$$\Gamma_{A_2} : \begin{array}{c} s \\ \bullet \end{array} \longrightarrow \begin{array}{c} t \\ \bullet \end{array} \quad A = \{s, sts\} \subseteq \{s, sts, t\} = T$$

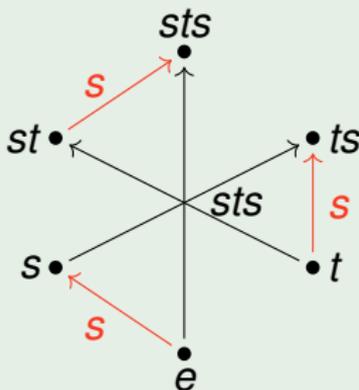


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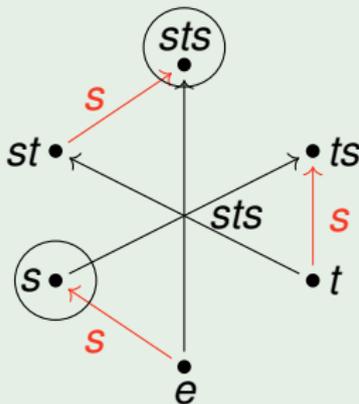


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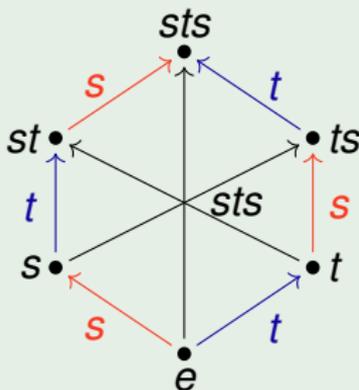


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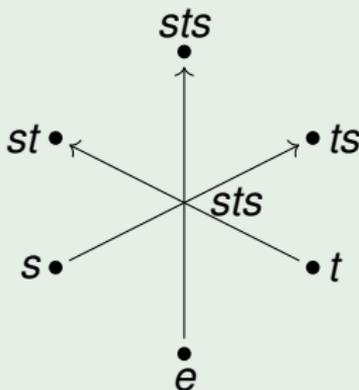


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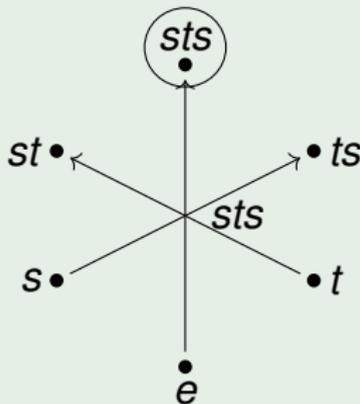


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## Bruhat preclosure

For  $A \subseteq T$

$$\bar{A} = \{t \in T \mid \text{there exists an } A\text{-path from } e \text{ to } t\}$$

### Example

$$\Gamma_{A_2} : \begin{array}{c} s \quad t \\ \bullet \text{---} \bullet \end{array}$$

$$\overline{\{s, sts\}} = \{s, sts\}$$

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$$\overline{\{s, sts\}} = \{s, sts\} = N(st)$$

$$\overline{\{sts\}} = \{sts\}$$

$$\overline{\{s, t\}} = \{s, sts, t\} = T = N(sts) = N(s \vee_R t)$$

## Dyer's conjecture

For  $A \subseteq T$  and  $u, v \in W$ :

1.  $\bar{A}$  is a closure operator.
2.  $\overline{N(u) \cup N(v)} = N(u \vee_R v)$ .

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## Dyer's conjecture

For  $A \subseteq T$  and  $u, v \in W$ :

1.  ~~$\bar{A}$  is a closure operator.~~ Can we do something about this?
2.  $\overline{N(u) \cup N(v)} = N(u \vee_R v)$ .

## Reflection order

A *reflection order*  $(T, \preceq)$  such that for every  $r, s \in T$ ,  
 $W' = \langle r, s \rangle$ :

$$r \prec rsr \prec \dots \prec srs \prec s$$

or

$$s \prec srs \prec \dots \prec rsr \prec r$$

### Example

$$\Gamma_{A_3} : \begin{array}{ccc} r & s & t \\ \bullet & \bullet & \bullet \\ | & | & | \\ \hline \end{array}$$

$$r \preceq t \preceq rstsr \preceq sts \preceq rsr \preceq s$$

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### Example

$$\Gamma_{A_3} : \bullet \xrightarrow{r} \bullet \xrightarrow{s} \bullet \xrightarrow{t}$$

$$\langle s, t \rangle \Rightarrow \{s, t, sts\}$$

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### Example

$$\Gamma_{A_3} : \begin{array}{ccc} r & s & t \\ \bullet & \bullet & \bullet \\ \hline & & \end{array}$$

$$\langle rsr, t \rangle \Rightarrow \{rsr, t, rstsr\}$$

$$r \preceq t \preceq rstsr \preceq sts \preceq rsr \preceq s$$

$$r \preceq rsr \preceq rstsr \preceq s \preceq sts \preceq t$$

## Initial section

Subset of  $T$  containing an "initial part" of \*some\* reflection order.

Finite = Inversion sets

### Example

$$\Gamma_{A_3} : \begin{array}{c} r \quad s \quad t \\ \bullet \text{---} \bullet \text{---} \bullet \end{array}$$

$$r \preceq t \preceq rstsr \preceq sts \preceq rsr \preceq s$$

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$\{r, t\}$	$\{r, rsr, rstsr\}$	$\{r, t, rstsr, sts\}$	$\emptyset$	$T$
$N(rt)$	$N(rst)$	$N(rtsr)$	$N(e)$	$N(rstrst)$

## Non-finite initial section: $\tilde{A}_1$

Subset of  $T$  containing an "initial part" of \*some\* reflection order.

### Example

$$\Gamma_{\tilde{A}_1} : \underset{\bullet}{s} \xrightarrow{\infty} \underset{\bullet}{t}$$

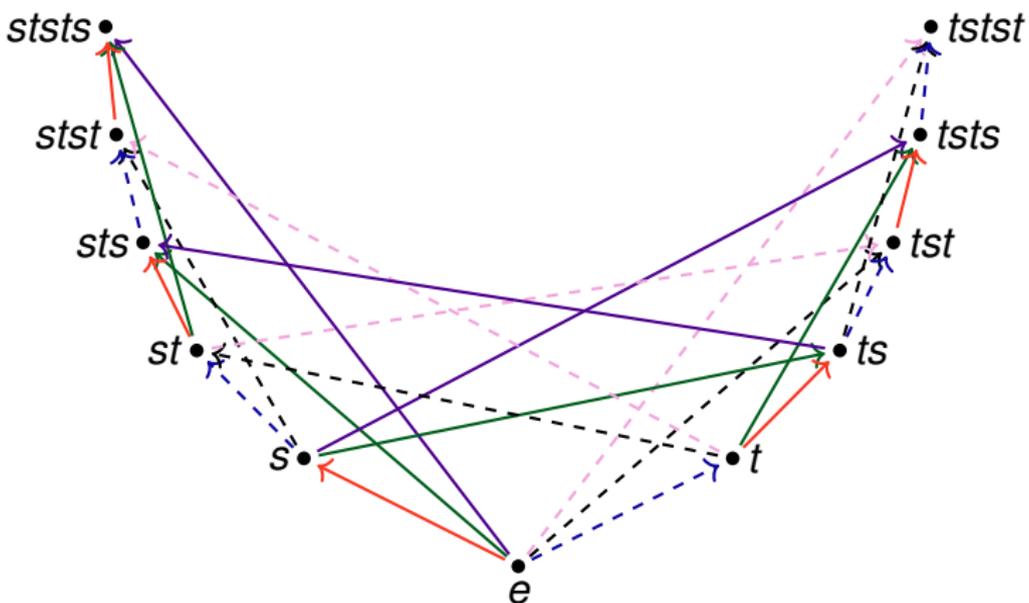
$$s \prec sts \prec ststs \prec \dots \prec tstst \prec tst \prec t$$

Initial section:  $\{w \in T \mid w \text{ begins with } s\}$

# Bruhat Preclosure

$$\tilde{A}_1 \quad \Gamma_{\tilde{A}_1} : \begin{array}{c} s \quad \infty \quad t \\ \bullet \text{---} \bullet \end{array}$$

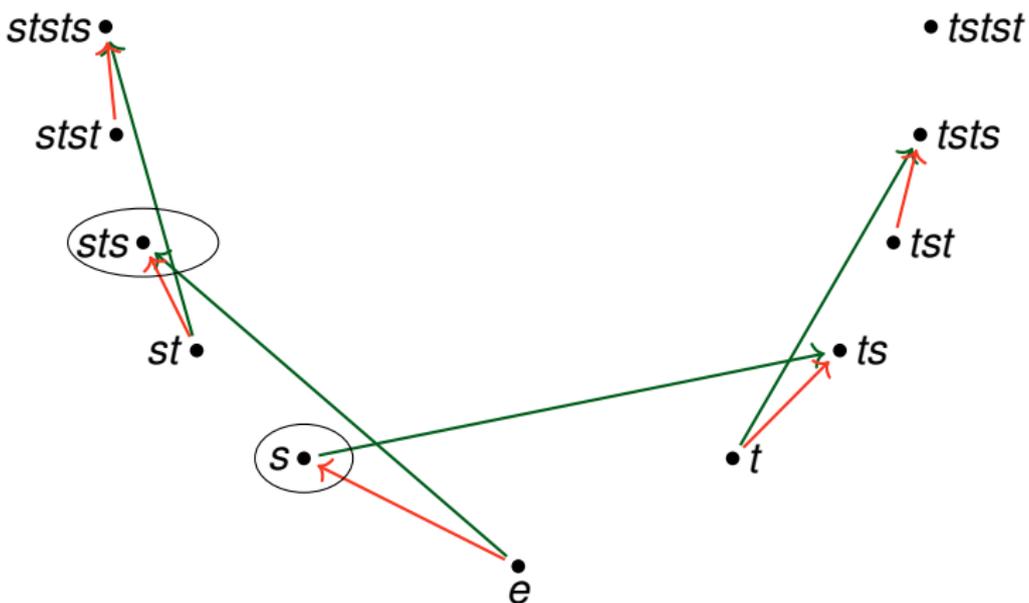
$s$  : red     $t$  : blue     $sts$  : green     $tst$  : black     $ststs$  : purple     $tstst$  : pink



# Bruhat Preclosure

$$\tilde{A}_1 \quad \Gamma_{\tilde{A}_1} : \begin{array}{c} s \quad \infty \quad t \\ \bullet \text{---} \bullet \end{array} \quad N(st) = \{s, sts\}$$

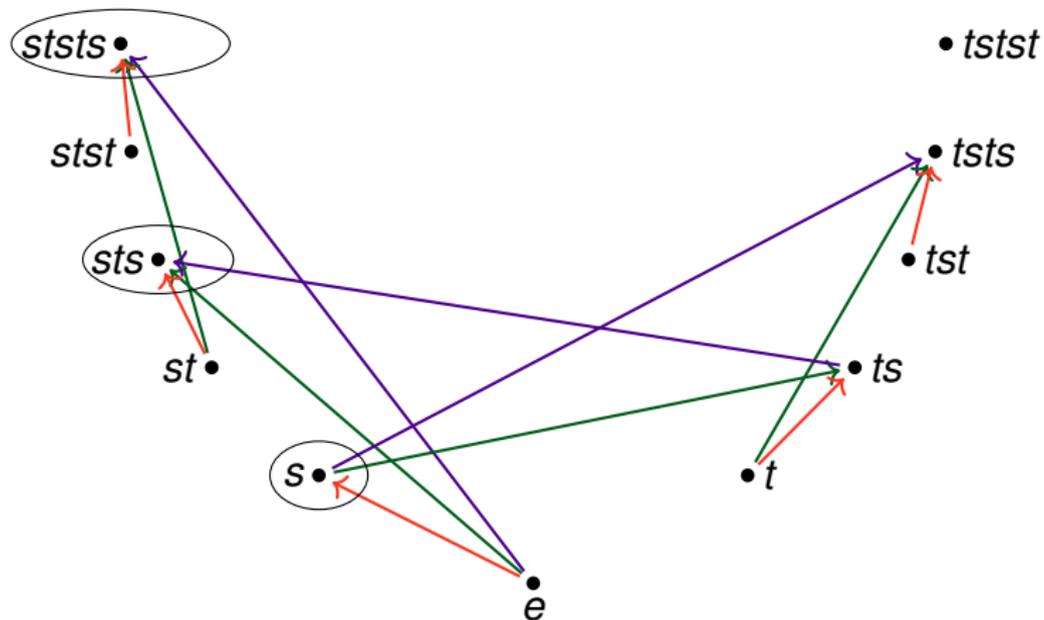
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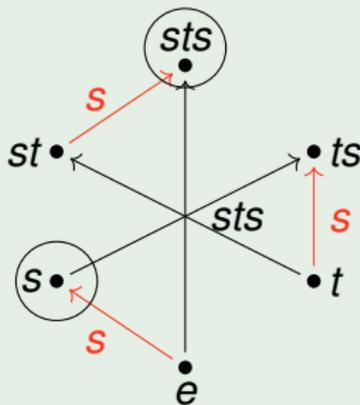
## Initial sections

Theorem (D. 2025+)

If  $A \subseteq T$  is an initial section, then  $\overline{A} = A$ .

### Example

$\Gamma_{A_2} : \begin{matrix} s & \longrightarrow & t \\ \bullet & & \bullet \end{matrix} \quad A = N(st) = \{s, sts\} \subseteq \{s, sts, t\} = T.$

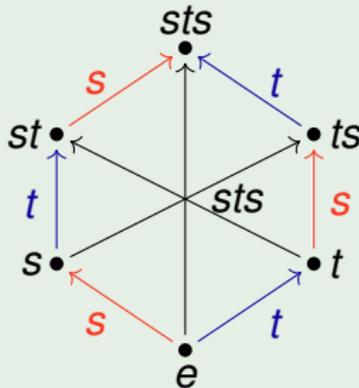


## Twisted Bruhat

For  $A \subseteq T$ , flip the edges labelled with elements of  $A$ .

### Example

$$\Gamma_{A_2} : \begin{array}{c} s \\ \bullet \end{array} \xrightarrow{t} \begin{array}{c} t \\ \bullet \end{array} \quad N(st) = \{s, sts\} \subseteq \{s, sts, t\} = T.$$

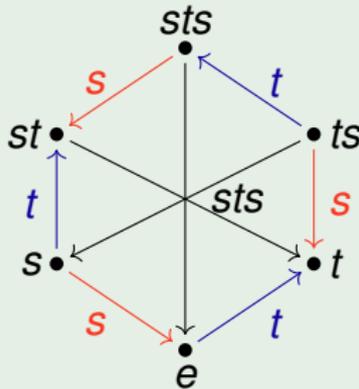


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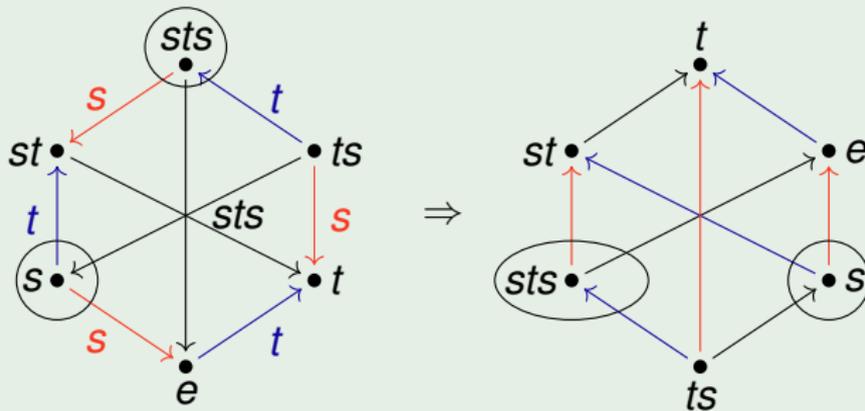


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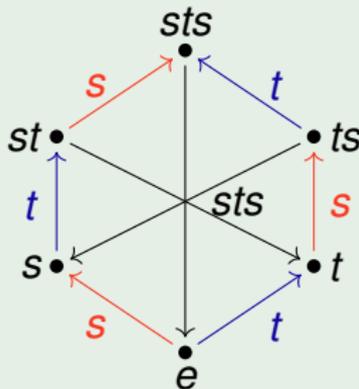
## Twisted Bruhat order

Theorem (Edgar '09 )

*The twisted Bruhat graph is a partial order for  $A \subseteq T$  if and only if  $A$  is an initial section.*

Example (Non-example)

$$\Gamma_{A_2} : \begin{array}{c} s \quad t \\ \bullet \quad \bullet \\ \hline \end{array} A = \{sts\} \subseteq \{s, sts, t\} = T.$$



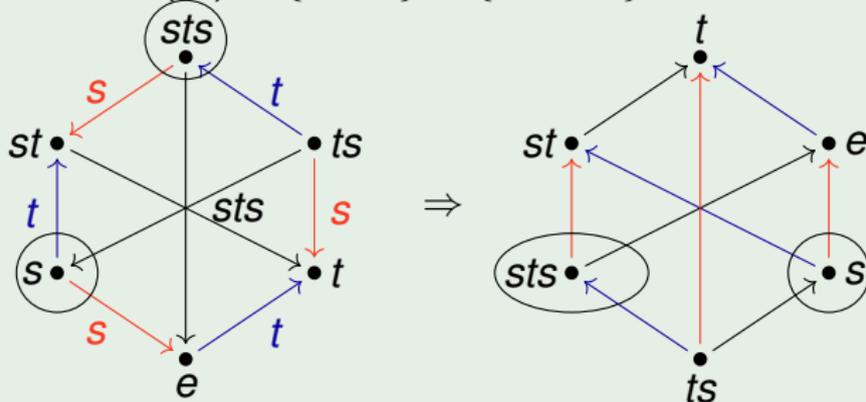
## Twisted Bruhat order

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### Example

$$\Gamma_{A_2} : \begin{array}{c} s \quad t \\ \bullet \quad \bullet \\ \hline \end{array} \quad N(st) = \{s, sts\} \subseteq \{s, sts, t\} = T.$$



## Twisted Bruhat order

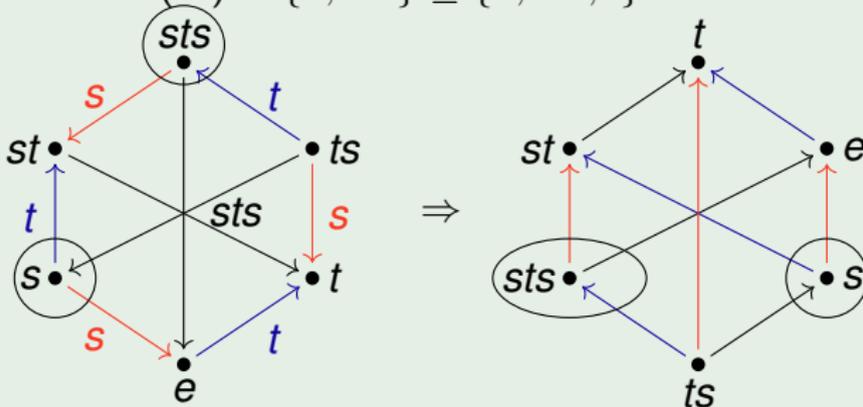
Theorem (D. 2025+)

For  $t \in T$  and  $A$  an initial section:

$$t \leq_A e \iff t \in \bar{A}$$

### Example

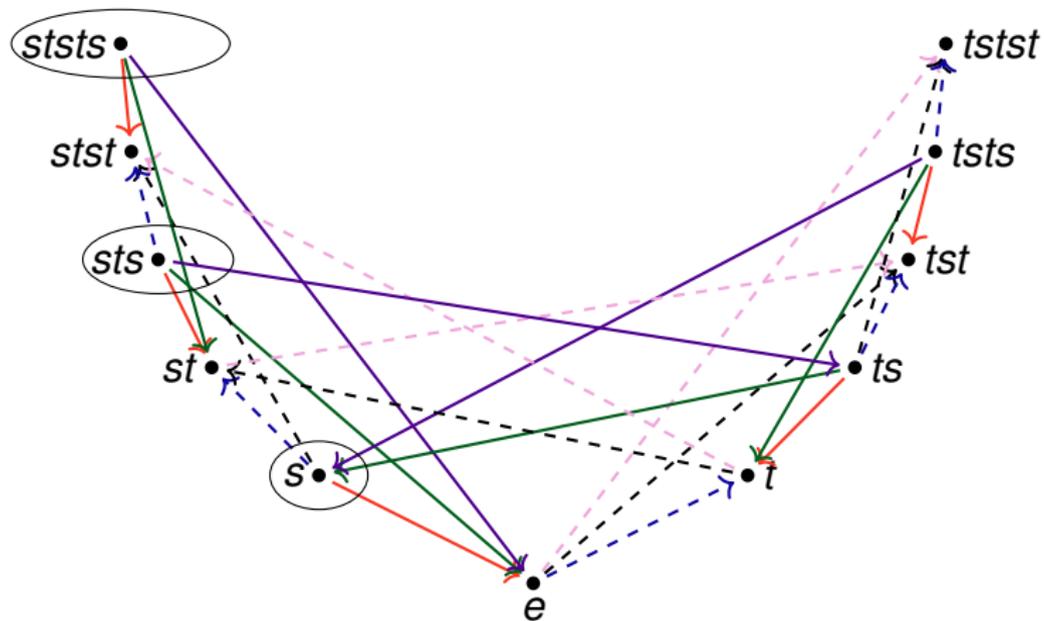
$\Gamma_{A_2} : \begin{array}{c} s \quad t \\ \bullet \quad \bullet \\ \hline \end{array} \quad N(st) = \{s, sts\} \subseteq \{s, sts, t\} = T.$



# Bruhat Preclosure

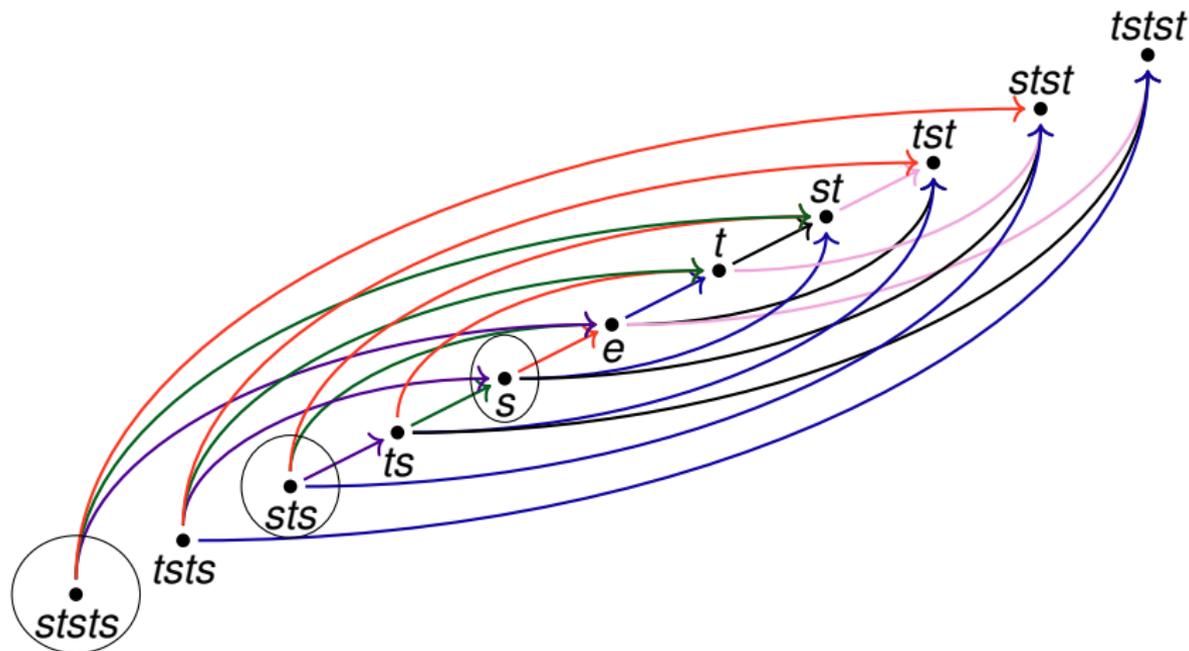
$$\tilde{A}_1 \quad \Gamma_{\tilde{A}_1} : \begin{array}{c} s \quad \infty \quad t \\ \bullet \quad \longrightarrow \quad \bullet \end{array} \quad \{w \in T \mid w \text{ begins with } s\}$$

$s$  : red     $t$  : blue     $sts$  : green     $tst$  : black     $ststs$  : purple     $tstst$  : pink



# Bruhat Preclosure

$$\tilde{A}_1 \quad \Gamma_{\tilde{A}_1} : \begin{array}{c} s \quad \infty \quad t \\ \bullet \text{---} \bullet \end{array} \quad \{w \in T \mid w \text{ begins with } s\}$$



## Bruhat order length

$$\ell_A(w) = \ell(w) - 2|N(w^{-1}) \cap A|.$$

## Example

$\Gamma_{A_2} : \begin{array}{c} s \quad \longrightarrow \quad t \\ \bullet \quad \quad \quad \bullet \end{array}$  For  $w = ts$  and  $A = \{s, sts\} = N(st)$  we have

$$\begin{aligned} \ell_A(ts) &= \ell(ts) - 2|N(st) \cap A| \\ &= 2 - 2|\{s, sts\}| \\ &= 2 - 4 = -2 \end{aligned}$$

## Theorem (D. 2025+)

For  $t \in T$  and  $A$  an initial section:

$$\ell_A(t) \leq 0 \iff t \leq_A e$$

## Initial section proof

Theorem (D. 2025+)

*If  $A \subseteq T$  is an initial section, then  $\bar{A} = A$ .*

$$\begin{aligned} t \in A &\iff \ell_A(t) < 0 && \text{(Dyer 1992)} \\ &\iff t \leq_A e \\ &\iff t \in \bar{A} \end{aligned}$$

## Finite type preclosure

$Z$  a set,  $\text{cl} : \mathcal{P}(Z) \rightarrow \mathcal{P}(Z)$  a preclosure.

$\text{cl}$  is of *finite type* if

$$\text{cl}(X) = \bigcup \{ \text{cl}(Y) \mid Y \subseteq X, |Y| \in \mathbb{N} \}$$

**Note:**  $\text{cl}^\infty(X) = \bigcup_{i \in \mathbb{N}} \text{cl}^i(X)$  is a closure.

Theorem (D. 2025+)

*Bruhat preclosure is of finite type.*

**Note:**  $\overline{A}^i$  is Bruhat preclosure performed  $i$  times:

$$\overline{A}^2 = \overline{\overline{A}} \quad \overline{A}^3 = \overline{\overline{\overline{A}}}$$

Corollary

For *any*  $A \subseteq T$ ,  $\overline{A}^\infty$  is a closure.

## Dyer's conjecture

For  $A \subseteq T$  and  $u, v \in W$ :

1.  ~~$\bar{A}$  is a closure operator.~~ Kinda fixed - Initial sections + Infinite Bruhat closure
2.  $\overline{N(u) \cup N(v)} = N(u \vee_R v)$ .

## Main theorem

## Theorem (D. 2025+)

Let  $W$  be an arbitrary Coxeter group. For  $u, v \in W$  such that  $u \vee_R v$  exists, then

$$\overline{N(u) \cup N(v)}^\infty = N(u \vee_R v)$$

## Corollary (D. 2025+)

For  $W$  a finite Coxeter group, then for all  $u, v \in W$

$$\overline{N(u) \cup N(v)}^i = N(u \vee_R v)$$

where  $i \leq |T \setminus (N(u) \cup N(v))|$ .

## Progress - Dihedral

$$\Gamma_{I_2(m)} : \begin{array}{c} s \quad m \quad t \\ \bullet \text{---} \bullet \end{array}$$

Theorem (Biagioli, Perrone 2025+, D. 2025+)

For  $W$  a dihedral group, for all  $u, v \in W$ ,

$$\overline{N(u) \cup N(v)} = N(u \vee_R v).$$

Progress - Symmetric group - type  $A_n$ 

Theorem (Biagioli, Perrone 2025+)

For  $W$  a type  $A_n$  Coxeter group, for all  $u, v \in W$

$$\overline{N(u) \cup N(v)} = N(u \vee_R v).$$

Theorem (D. 2025+)

For  $W$  a type  $A_n$  Coxeter group and  $A \subseteq T$ ,  $\bar{A}$  is a closure.

## Progress - Other types

Finite types:

$B_n$  Recently proved by Biagioli, Perrone 2026+

$D_n$  **Still open** (Checked for  $n \leq 5$ )

$F_4$  Computer check by Biagioli, Perrone 2026+

$H_3$  Computer check by Biagioli, Perrone 2026+

$H_4$  Computer check by Dermenjian 2026+

$E$  **Still open** (Check for  $E_6$  in progress)

Infinite types:

- **Everything open**

Thank you

